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(11) EP 1 022 860 A1

(12)

EUROPEAN PATENT APPLICATION

(43) Date of publication: 26.07.2000 Bulletin 2000/30

(21) Application number: 00101174.1

(22) Date of filing 21.01.2000

(51) Int. CI.7: **H03M 13/29**

This Page Siank (uspio)

(84) Designated Contracting States:

AT BEICH CY DE DKIES FIFR GBIGRIE IT LILLU MC NL PTISE
Designated Extension States

ALLT LY MIX ROSI

(30) Priority 22 01 1999 US 116765 P 28.07 1999 US 363303

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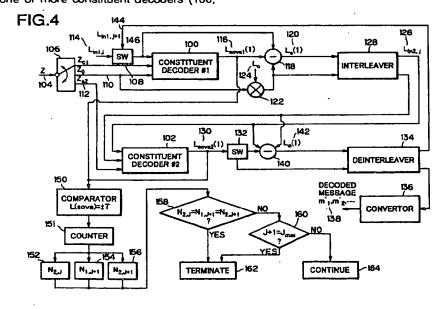
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(54) Turbo decoding with variable number of iterations

(57) Turbo codes used in digital communication systems are decoded by successive iterations of a maximum likelihood algorithm. The turbo code decoder is a modular apparatus with each module comprising two or more serial constituent decoders (100, 102). The output a constituent decoder (100) or a module is part of the input to the next constituent decoder (102) or module, respectively. The progress of the decoding can be monitored by establishing a limit for the output of the constituent decoders (100, 102) and monitoring the numbers of outputs of one or more constituent decoders (100,

102) that approximately equaling that limit. When the numbers of outputs approaching the limit or saturating is unchanged for successive iterations or successive serial decoders (100, 102) no further progress is being made in decoding the message. The performance of the turbo code decoder can be optimized by terminating iterations of the decoding process when the numbers of constituent decoder outputs that have saturated is unchanged for successive iterations or decoder operations.



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BACKGROUND OF THE INVENTION

The present invention relates to channel coding for digital communications systems and, more particularly, to a turbo, code decoder useful for channel coding and a method for optimizing the performance of a turbo code ya <mark>degoder</mark>a ji bili beki ingusiran ingusira ya manara sababbi birka bira tahar bir bir kimula masa sa ingusira y [0902] Forward error correction (FEC) is a system of error control for data transmission systems where the receiver and is capable of detecting and correcting errors in the "astreceived" message induced by noise in the transmission chan-10, and EEC is useful in connection with data transmission systems which lack a reverse channel with which retransmission of data can be requested or where retransmission would be difficult because the delay would be excessive or repeated a retransmission would be required because of the number of expected errors. For these reasons, FEC has been of parestaticular interest and use in wireless communication and space probe and satellite data transmission systems. FEC relies channel coding where input message sequences are mapped to code symbol sequences that add redundancy and memory to the data before transmission. Connected to a source of the second contract 3. [0003] Conference Generally, channel coding utilizes block or convolutional coding to add redundancy to the message bit gestream. Block coding breaks the bit stream representing the message into fixed size blocks and independently adds the gradundant code symbols to each block. Block coding is usually decoded with algebraic techniques. On the other hand, 🚋 convolutional coding continuously adds redundant symbols to the bit stream based on the contents of the stream: In the 20 geconvolutional encoder, the bits of the message are shifted serially into and out of a shift register having a number of , individual registers. The output code is the result of modulo; arithmetic performed on the contents of the shift register and, in some cases, the input bit stream as each successive message symbol or bit is shifted into the register. While bit stream segmentation is not required for convolutional coding the coded bit stream is typically broken into blocks or restrames for other reasons before transmission: Discoding of convolutional codes is accomplished with a heuristic or trial-25 wand-error approach in the response to the interest tests after your ensisted as a treatment to the interest tests and the artificial and the a [0004] Turbo codes are produced by encoders comprising two, or more, parallel, constituent encoders. The constituent uent encoders are often, but not necessarily, identical convolutional encoders. An interleaver or permuter is attached to the input of one or more of the constituent encoders. The interleaver rearranges the input of the attached constituent ecencoder in a systematic, pseudo-random manner. As a result, turbo codes comprise two or more (depending on the 30 sampler of encoders); independently coded symbol streams that refer to the same input information bit stream Turbo an society of particular interest because with a relatively simple constituent code and large interleaver their performance encap benearathe theoretical or Shandon limit of the transmission channels abby \$1.00 generates native whose byte 44 [0005] 70 Turbo code decoding is an iterative process with the results of a first modular decoder forming part of the input to a second modular decoder and so forth until the required number of iterations is achieved. When the turbo code 35 is composed of two parallel concatenated codes the modular turbo code decoder comprises two serially conflected constituent decoders separated by an interleaver that reorders the output of the first decoder so that it may be used as input to the next decoder. Decaders for turboccodes with more than two parallel constituent codes may take a number and forms. A convolutional encoder is a state machine that codes by tracing a path through a code tree or trellis on the basis of the sequence of input symbols. From the symbols of the "as received;" coded message the donvolutional code 40 predecoder attempts to retrace the encoder's path through the code tree or trellis outputting the symbols of the decoded message while correcting errors incurred in transmission. One technique for decoding convolutional codes relies on algorithms which retrace the path of "maximum likelihood" through the trellis. One such "maximum likelihood" algorithm sused in turbo code decoding is the soft output Viterbi algorithm (SOVA). A constituent decoder applying the SOVA algorithm computes or estimates the "log likelihood ratio," the logarithm of the ratio of the conditional probabilities of receiv-45 sing the two outcomes (binary 11" and "0") given the observed signal value. The output of the constituent decoder is a suplurality of signed numbers. The sign expresses the polarity of the decoded message symbol. The magnitude is a "soft" or analog value expressing the probability that the decoded symbol is the same as the original message symbol. Generally, the turbo code decoder converges on a final decoded symbol sequence with successive iterations and the error rate performance improves until a threshold number of iterations is reached. While the error rate performance of the decoder generally improves with additional iterations, the rate of improvement decreases. Each iteration takes time further delaying completion of decoding. Heretofore, the number of iterations to be performed by a particular turbo code decoder was hard wired into the decoder. Optimizing the number of iterations to be hardwired into the decoder involves compromises in the error rate and latency of the decoder's performance. Further, due to the random nature of noise, "as received":data sequences are unequally corrupted and require different numbers of iterations 55 to achieve the same level of error correction. The works were the same of action control is the control of t What is desired therefore; is a turbo code decoder and a method of decoding that optimize the performance of the decoder producing a given level of error correction in the fewest number of decoding iterations on the average.

Further, it is desired that the operation of the decoder be responsive to the error correcting requirements of each mes-

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sage. Optimizing the decoding process reduces the latency in decoding a message at an acceptable error rate and reduces the cost and complexity of the decoder.

SUMMARY OF THE INVENTION

មាននៃខាត់ តា តាមប្រភព្ធមាន ខាត់ មិន សាមការប្រកាសន៍នៃ មានប្រកាសន៍ and the experience of the experience of [0008] :The present invention overcomes the aforementioned drawbacks of the prior art by providing a method of optimizing the performance of an iterating turbo code decoder having a plurality of constituent decoders comprising the steps of establishing a limit for an output of a constituent decoder, determining the number of decoder outputs approximately equaling the limit for each of successive iterations by the turbo code decoder, and terminating operation of the turbo code decoder when the numbers of the decoder outputs approximately equaling the limit are substantially - unchanged for successive iterations. The progress of a turbo code decoder in decoding a message sequence can be remonitored by establishing a limit for the output of a constituent decoder and monitoring the number of outputs of equaling or approaching the limit. If The number of outputs approaching the limit or saturating does not change for successive by literations by the turbo code decoder inciprogress is being made in decoding the message. Operation of the turbo code decoder can be terminated without loss of data. 38. [0009] A second technique applying the method of optimizing the performance of an iterating turbo code decoder ¿ including a plurality of serial constituent decoders comprises establishing a limit for an output of a constituent decoder a determining a first, and a second number of the outputs approximately equaling the fimit produced by the second serial constituent decoder while performing a first iteration and a second iteration, respectively; determining a third number of 20 the outputs approximately equaling the limit produced by the first serial constituent decoder while performing the secet and iteration; and terminating operation of the turbo code decoder when the first, second, and third numbers of outputs e y l<mark>are substantially/equal</mark>vernita at lid to lod ova egulta an er/esecoule i noe i e. . . . 55 .[6010] 201 Arthird technique for applying the method of optimizing the performance of an iterating turbo code decoder mincluding a plurality of serial constituent decoders comprises establishing a fimit for an output of a constituent decoder: determining a first number of outputs produced by the first serial constituent decoder performing an iteration approximately equaling the limit; determining a second number of outputs produced by the last serial constituent decoder performing the iteration approximately equaling the limit, and terminating operation of the turbo code decoder when the first Chapter and the first of the second to the s and second numbers of outputs are substantially equal. [001/1] one local fourth technique of applying the method of optimizing the performance of an iterating turbo code ുപ്പം decoder, including a plurality of serial constituent decoders comprises establishing a limit for an output of a constituent so decoder, and terminating operation of the turbo-code decoder when substantially all outputs produced by the constituent decoder while performing an iteration are approximately equal to the limit. By monitoring the progress of the decodeming process and terminating decoding iterations when progress slows or stops, the delay in decoding any particular comessage stream can be minimized while achieving a satisfactory error rate for any "as received" symbol sequence. 35jer [0012]: An optimized iterating turbo:code decoder/is also provided comprising a plurality of serial constituent decoder ge ters; a comparatento compare a plurality of outputs of a constituent decoder to a threshold value for the output, a counter souto, count the number of the outputs produced by the constituent decoder approximately equaling the threshold Value; get and a decision unit to terminate operation of the turbo code decoder when the number of the outputs approximately att/equaling the threshold value is, unchanged for successive iterations. Int a second embodiment of the optimized turbo 40,ecode decoder, the number of outputs of the first and last of the serial constituent decoders that have saturated during no agriteration is compared to determine progress in decoding, final third embodiment, the number of outsuts of the first serial constituent decoder performing a first iteration and the numbers of outputs of the last serial constituent decoder empedorming the first and a subsequent iteration are compared to determine decoding progress and whether further itervisiations are warranted robbiness is into a time of the methispoliente, of the body of the assembles is autouprometric 45; g-[00]3]shim: The foregoing and other objectives; features and advantages of the invention will be more readily underenstood upon consideration of the following detailed description of the invention, taken in conjunction with the accompanying drawings, early benefind any an amballed religious general to the solid precision of the case of the artists of the solid precision of the case Que et dy fina dama con a cooder, converant no cinari, tenocud ev<mark>rantol secucinos cum</mark>a que rissi e raje g (BRIEF DESCRIPTION: OF: THE DRAWINGS \$ 0.5 or the covered name of common sections of the sections and the arms 50%. 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and FIG. 2to arista block-diagram of an exemplary systematic; recursive convolutional encoder. All soften in which is often a contract.

്യെFIG. 5-നു is trellis_diagram_for a soft-output Viterbi algorithm-based, constituenty convolutional decoderant വാധം പ്രസംഗുള്ള നിഴുന്നുണ്ട് വെന്നുണ്ടെ വരുന്നു വരുന്നു. പ്രവേശന പ്രശ്നായായില് വരുന്നു വരുന്നു വരുന്നു വരുന്നു വരുന

FIG.4 acrosis: a block diagram of a turbo; code decoder according to the present invention (1986) 1997 and

is a trellis diagram of the operation of a recursive, constituent convolutional encoder of the encoder of FIG. 2.

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TOODETAILED DESCRIPTION OF THE PREFERRED-EMBODIMENT (DOGGLES) OF THE SECOND ROOM

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[0016] Several channel coding techniques are used for FEC in communication systems. One technique, which is thought to achieve performance near the theoretical of Shannon limit of the transmission channel, is turbo coding. Turbo codes are interleaved, parallel concatenated, codes. The constituent codes are offen convolutional codes but may be codes produced by other techniques.

[0017] Referring to FIG. 2, an exemplary turbo-code encoder 40 (indicated by a bracket) comprises two parallel, recursive, systematic, convolutional, constituent encoders 42 and 44 (indicated by forackets) with an interleaver 46 attached to the input of one of the encoders 44. Turbo code encoders may have more than two parallel constituent encoders with additional interleaver attached to the inputs of the additional constituent encoders with additional interleaver attached to the inputs of the additional constituent encoders. The turbo code encoder 40 has three intermediate outputs (c₀) 47, (c₀) 48, and (c₀) 50 which are combined into a coded serial bit estream (C) 52 by an output multiplexer-54. The code rate of the turbo code encoder 40 of the ratio of the number of input in elessage symbols (m) to the number of output code symbols (C) is 1/3. The code symbols (C) i

The turbo code encoder 40 is a systematic encoder with the symbols of the input message of frame (m) 56 constituting one of encoders intermediate outputs (c₀) 47. The intermediate outputs (c₀) 48 and (c₀) 50 are the outputs of the first 42 and second 44 constituent convolutional encoders, respectively. The first 42 and second 44 constituent convolutional encoders as shift register 58 and a pair of modulo-2 adders 60 and 62. The shift registers 58 of the illustrated encoders have two individual registers; (T1) 64 and (T2) 66. The symbols of an input bit stream for the constituent encoders 42 and 44, (m) 56 or (m(a)) 68, respectively, are shifted one bit at a time into the shift register 58. To ensure that the conferts of the shift register 58 are known at the start of the message and that all of the message bits are shifted completely through each shift register 58, the turbo code encoder includes a padding unit-70 to add sufficient bits (usually zeroes) to "flush" the shift register to a known state at the end of the message or frame. A convolutional encoder adds redundant bits to the message on a continuous basis so that the entire message as a sone code. However, the coded message data is often segmented into fixed length frames for transmission.

[0019] In the turbo code encoder 40 the input of the first constituent convolutional encoder 42 is the exemplary message sequence (m) 56[10101] In the recursive convolutional encoders 42 and 44 the modulo-2 sum of the input bit and the contents of the two registers (T1) 64 and (T2) 66 is computed in a modulo-2 adder 60 to create a feedback bit 72 (F). The output of the constituent convolutional encoder, (c_{C1}) 48 or (c_{C2}) 50, is the modulo-2 sum of the feedback bit 40 972 (F) and the contents of the T2 register 66 computed in the adder 62. The output of the first constituent convolutional encoder (c_{C1}) 42 is the coded bit stream: [11011].

[0020] The input (m(a)) 68 of the second constituent convolutional encoder 44 is the message sequence (m) 56 as rearranged by an interleaver 46. Interleaving rearranges the input symbols in some pseudo-random manner so that the seminated from each other in the output bit stream (C) 450-52. The error correcting performance turbo codes is due, in part, to the use of large turbo interleavers that interleave hundreds bits interleaving results in two or more (depending on the number of constituent codes) independently coded symbol streams that refer to the same message bit stream. Turbo code decoders exploit this structure of the code to message bit stream from the bits in the remaining stream(s) and use this information to estimate the correctly decoded symbols. The estimates are refined during successive iterations of the decodes.

[0021] In the exemplary interleaver 46 the input symbols are rearranged according to a table where the first bit of the input sequence (m₁) becomes the fourth bit (m(a)₄) of the convolutional encoder input sequence 68; the second input bit (m₂) becomes the third bit (m(a)₃) and so forth. As a result the output (c_{c2}) 50 of the second convolutional encoder 44 is different from that of the first constituent encoder 42 even though the encoding process within the second convolutional encoder is the same. Further as a result of interleaving, the three coded output symbols representing any input message symbol are separated in the output bit stream 52.

[0022] A convolutional encoder is a state machine and its operation can be represented by a path traced through a code free or a trellis diagram. FIG. 3 is a trellis diagram illustrating an exemplary operation of the constituent recursive

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convolutional encoder 42 of FIG. 2. Each node 80 of the trellis corresponds to a state of the shift register of the encoder 42 (listed vertically to the left of the trellis) at an input time or coding stage 82 (to, t1,...) indicated horizontally across the top of the trellis). Two paths exit each node of the trellis. The path to be followed to the next state at the next stage of encoding is determined by the message bit causing the state transition. An input message bit (1) causes the encoder to follow the path represented by a dashed line and an input message bit (0) causes the encoder to follow the path represented by a solid line. An output code symbol 84 for each path is shown adjacent to the path. For example, at the bit (1) is shifted into the encoder 42 causing the encoder to transition from the an initial state (00) 86 (as a result of padding) to the next state (10) 88. When the encoder exits the first node (00) 86 for the next state (10) 88, it exits on the path corresponding to the input bit (1) (dashed line) which causes the encoder to output a code symbol (11) 90 Likewise, the next message bit (0) determines the exit path from the second node (state 10 at t1) 88 producing the next code symbol , (01) 92 and causing the encoder to move to its next state (11 at t2) 94, It is desirable for the encoder to start and end in an all zero state. For an input sequence [101] two extra bits [01] are necessary to produce a zero ending state. Following the path through the trellis for an input message sequence [101] plus two padding bits [0,1], produces the output code: [1101100111]. The output of the encoder of FIG. 2 is the output constituent encoders (the trellis diagrams for 15... each of the two encoders) which is multiplexed with the original message symbols. definitions in

Jurbo code decoding is an iterative process and the turbo code decoder is modular in nature with the output of a first module (the result of a first iteration) being the part of the input to a first decoder of the next modular decoder for use in the next iteration of the process. Referring to E.G. 4, an exemplary modular turbo code decoder includes a number of serially connected, constituent decoders 100 and 102 equal to the number constituent encoders in the particular turbo code encoder. Other arrangements of the decoder are possible if the number of constituent codes is greater than two, but the present invention may be utilized with these other arrangements, as well. The "as received" coded transmission (Z) 104, including any errors incurred in transmission, is received by the decoder from the demodulator, 22, lip, an input demultiplexes 106, the "as received" catta stream. (Z) 104 is separated into three constituent streams; (z₀) 110, (z₀) 108, and (z₀) 50 dt is a received as received versions of the three intermediate outputs of the encoder (c₀) 47, (c₀) 48, and (c₀) 50 dt is a received as received.

[0024] In the constituent decoders 100 and 102, the soft output Viterbi algorithm (SOVA) is used to retrace the "most likely" path followed by the encoder through the trellis diagram when it produced the encoded message. Referring to FIG.5, in applying the Viterbi algorithm the decoder operates serially through the stages, At each stage, a path metric 170 is computed for paths from the initial state to all possible states for that stage. The path metric 170 is the accumulated difference metric, 172 along each path from the decoder's initial state 174 to each node 176 at the particular stage. The difference metric, 172 for a path between two nodes, represents a measure of the difference between a code word of the received message and the corresponding code word which the encoder would have generated in making the stransition between those two nodes or states. According to the Viterbi algorithm, the "most likely" path followed by the encoder in coding the sequence is the path with the least path metric. This path is retained for use at the next input stage and the path with the greater path metric is discarded. As the algorithm is applied for input stages further and further into the trellis a single "most likely" path emerges after a delay 180. There is a high probability that errors in the path the diagram ago. 103 along the stage and that the message will be correctly decoded by following this path through the strellis diagram.

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[0025] ___A_modified_Viterbi algorithm (the soft-output Viterbi algorithm or SOVA) is used in turbo code decoders to produce "soft" or analog outputs from soft or hard inputs. "Soft" inputs and outputs quantize the level of the actual signal (often at eight levels) between the possible a binary values of the signal. The "soft" value represents the reliability of the signal's value as well as its polarity.

[0026]. There are two basic steps in the SOVA algorithm used in the constituent decoders of the turbo, code decoder. The first or Viterbi step is similar to the Viterbi algorithm with the addition of the computation of the maximum path metasic differences for each node at each stage of the trellis. In the second or update step, a window 178 is established corresponding to the delay of the Viterbi algorithm in converging on a single path. The window 178 slides along the trellis as the decoder moves from stage to stage. The minimum path metric difference is found from among all possible paths in this window that compete with the survivor path and would have led to a different bit decision than the survivor path. This minimum path metric difference is given a sign based on the hard value of the current bit and is the soft value of the SOVA output The output of a constituent convolutional decoder applying the SOVA algorithm comprises a signed number for each decoded bit of the message. The sign of the number (the sign of the "log likelihood ratio") represents the value of the "log likelihood ratio") represents the reliability of the plikelihood ratio").

10027] 386 With the SOVA algorithm, the path metric computations of the Viterbi algorithm are augmented by the addition of a term which incorporates extrinsic information (Lin(t)) supplied by the previous constituent decoder. For each state (k) at a particular t in the trellis, the path metric M_{k,t} is computed by extending the path metrics from all states (k'), at time t-1, with valid transitions to a state k, at time t, as follows:

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20 [0028] Referring to FIG. 4, the first constituent decoder 100 applies the SOVA algorithm to the systematic data (Z₀)110, the coded data from the first encoder (Z_{C1}) 108, and initial a-priori information (L_{in1,i}) 114 to compute a soft output
(L(sova1) 116 comprising a plurality of signed numbers each representing the value of a decoded message bit and the
probability that the value has been correctly decoded. The extrinsic information (L(e)):120 is produced by subtracting
the product of systematic data (Z₀) 110 and the channel reliability factor (Lc) 124 produced in the multiplier 122 and the
a-priori information (L_{in1,i}) 114 from the decoder output (L(sova)):116 in a subtraction unit 118. After interleaving the
extrinsic information (L(e)) 120 becomes a priori information (Lin2,j):126 which is supplied to the second constituent
decoder 102.

[0029] The input data for the second constituent decoder 102 comprises the interleaved version of a priori information (Lin2,2j) 126 from the first constituent decoder 100, the systematic data (Z₀) 110 which has also been interleaved in the interleaver 128, and the "as received" coded output of the second encoder of the turbo code encoder (Z_{C2}) 112 which has been separated from the input message bit stream; (Z) 104 by the input demultiplexer 106. The second constituent decoder 102 applies the SOVA-algorithm to this second coded information and produces a second soft output (Lsova2) 130, completing an iteration of the decoding process.

[0030] gis lift the required number of iterations have been completed, the soft output of the second decoder (Lsova2)
130 is diverted by switch 132 to the deinterleaver 134 and passed to an conversion unit 136 where "hard" (binary) val-

[0031] If the required number of iterations has not been completed, the interleaved intrinsic information (Lin2.j) and the interleaved systematic data (Z₀) 110 from the input to the second decoder 102 is subtracted from the soft output of the second decoder (Lsova2):130 in the subtraction unit 140. The resulting extrinsic information (Lin 142 is an inter-

46. \leaved version of the intrinsic information (bin3)#1)*1445After deinterleaving in the deinterleaver 134, this intrinsic information:(bin3)#1)*144 is fed back through the switch 146 to the input to the first constituent decoder 100 for the next successive iteration of the decoding process.

[0032] Fall The output of the SOVA algorithm depends upon the value of the minimum of the maximum path metric differences in the update window along the survivor path. For low signal to noise ratios (SNR) the maximum path metric differences are likely to be small. For high SNR, the path metric differences are likely to be large resulting in large values of the SOVA output (Lsova). These large values have a negative effect on the convergence of the algorithm and the performance of the algorithm can be improved by limiting the output (Lsova).

[0033] we The present inventor came to unexpected realization that the convergence of the SOVA algorithm can be monitored by comparing the number of outputs of the constituent decoders that are saturating at a limiting level at each iteration and discontinuing iteration when this number does not change from iteration to iteration. Further, the decision to continue iterating dan be made by comparing the change in the number of saturating outputs produced by each of the individual constituent decoders in a single iteration.

[0034] What the present invention; the outputs (L(sova)) are directed from the constituent decoders 100 and 102 to a comparator 150 which compares each output to a threshold or limit (T). The numbers of L(sova) outputs from the second constituent decoder 102 (L(sova2)) for a first iteration (N2,j) and a successive iteration (N2,j+1) and from the first constituent decoder 100 for the successive iteration (N1,j) equaling the threshold (T) are counted by the counter 151 and stored in memory 152 (154) and 156. The three numbers of saturated L(sova) values are compared in a decision unit 158. If the three numbers are approximately equal; the algorithm can be declared to have converged and the oper-

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ation of the turbo code decoder can be terminated 162 without loss of information. If the three numbers are not equal, the algorithm is continuing to converge and the turbo code decoder will be allowed to continue 164 to operate unless the maximum number of iterations 160 has been reached. A second technique and apparatus stores and compares the numbers of outputs of the first and second constituent decoders which have saturated in successive iterations. When the numbers of saturating outputs do not in succeeding iteration does not change convergence is declared and operation of the decoder is halted. Likewise, the saturating outputs of a single constituent decoder can monitored for successive iterations until there is little or no further change in the numbers of saturating outputs. A fourth, less complex technique is to compare the output of the last serial constituent decoder with the output of the first serial constituent decoder performing one iteration. While the convergence of the algorithm is indicated by a lack of change in the number of saturated outputs produced by the first and last decoders performing a single iteration; there is less certainty of the effect of continuing operation using this method. If all of the values in the SOVA outputs of the constituent decoders have saturated the algorithm has converged and operation may be terminated. The threshold (T) and the L(sova) value are both signed values. The number of comparisons may be reduced by 50% by comparing outputs of only one sign with threshold having the same sign.

[0035] The terms and expressions that have been employed in the foregoing specification are used as terms of description and not of limitation, and there is no intention, in the use of such terms and expressions, of excluding equivalents of the features shown and described or portions thereofait being recognized that the scope of the invention is defined and limited only by the claims that follow.

- 20 (Claims strate, and or refringer AVOS and solding 20 hebroad requirements. 1997)

 the probability of the performance of an iterating turbor code decoder including a plurality of constituent grand addecoders (100, 102) comprising the performance of an iterating turbor code decoder including a plurality of constituent grand addecoders (100, 102) comprising the steps of this self. Inches by the analyse and present an order of an iterating turbor code decoder including a plurality of constituent grand addecoders (100, 102) comprising the steps of this self. Inches by the analyse of the
- The providers (100, 102) comprising the activation and sententies of the sententies

4. A method of optimizing the performance of an iterating turbo code decoder including a plurality of serial constituent

- in the substantially equal to the constituent of said outputs approximately equaling said limit produced by a first said serial states a constituent decoder (100) while performing said second iteration; and the second s
- Table A Hope will you be to come the condition of the con
- 55min of the Stabilishing and the performance of an iterating furboloode decoder including a plurality of serial constituent of serial constituent of serial constituent of serial constituent is an including a plurality of serial constituent of serial constituent of serial constituent decoder (1,00, 1,02); and the serial constituent of serial constituent decoder (1,00, 1,02); and the serial constituent decoder (1,00, 1,02); and the serial constituent decoder when substantially allisaid outputs produced by said constituent decoder (1,00, 1,02); and the serial constituent

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- The method of claim 7 wherein a said constituent decoder (100, 102) applies the soft output Viterbi algorithm.
 - 10. A method of optimizing the performance of an iterating turbo code decoder including a plurality of serial constituent decoders (100, 102) comprising:
 - (a) establishing a limit for an output of a said constituent decoder (100, 102);
 - (b) determining a first number of said outputs approximately equaling said limit produced by a first said serial constituent decoder (100) performing an iteration;
 - (c) determining a second number of said outputs approximately equaling said limit produced by a last said serial constituent decoder (102) performing said iteration; and
 - (d) terminating operation of said turbo code decoder when said first and second numbers of said outputs are substantially equal.
 - 11. The method of claim 10 wherein said first and said second numbers of said outputs are numbers of said outputs of one sign approximately equaling said limit having said same sign.
- 20 12. The method of claim 10 wherein a said constituent decoder (100, 102) applies the soft output Viterbi algorithm.
 - 13. An iterating turbo code decoder comprising:

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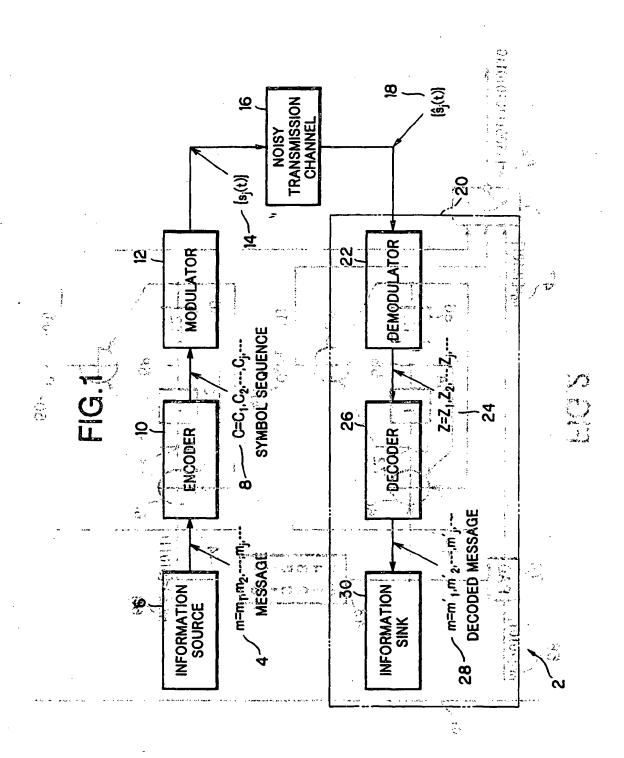
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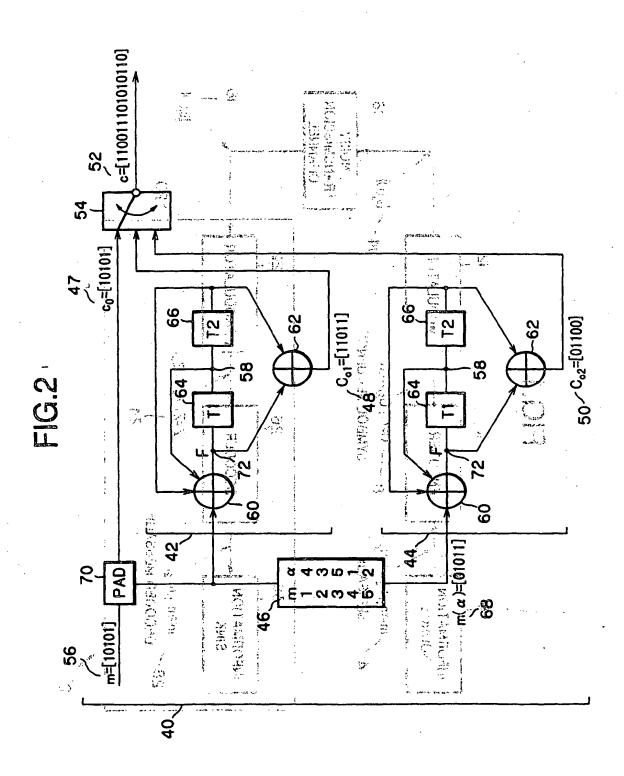
- (a) a plurality of serial constituent decoders (100, 102),
- (b) a comparator (150) to compare a plurality of outputs of a said constituent decoder (100, 102) to a threshold value for said outputs;
- (c) a counter (151) to count a number of said outputs produced by said constituent decoder (100, 102) approximately equaling said threshold value; and
- (d) a decision unit (158) to terminate operation of said turbo code decoder when said number of said outputs approximately equaling said threshold value is unchanged for successive iterations.
- 14. The turbo code decoder of claim 13 wherein said constituent decoder (100, 102) is a soft output Viterbi algorithm decoder.
- 35 15. An iterating turbo code decoder comprising:
 - (a) a plurality of serial constituent decoders (100, 102),
 - (b) a comparator (150) to compare an output of a said constituent decoder (100, 102) to a threshold value for said output,
 - (c) a counter (151) to count a first and a second number of said outputs produced, respectively, by a first said serial constituent decoder (100) and a subsequent said serial constituent decoder (102) approximately equaling said threshold value; and
 - (d) a decision unit (158) to terminate the operation of said turbo code decoder when said first number of said outputs is approximately equal to said second numbers of said outputs.
 - 16. The turbo code decoder of claim 15 wherein said constituent decoder (100, 102) is a soft output Viterbi algorithm decoder.
 - 17. An iterating turbo code decoder comprising:
 - (a) a plurality of serial constituent decoders (100, 102),
 - (b) a comparator (150) to compare an output of a said constituent decoder (100, 102) to a threshold value for said output.
 - (c) a counter to count a number of said outputs approximately equaling said threshold value;
 - (d) a memory (152, 154, 156) to store a first, a second, and a third said number of said outputs approximately equaling said threshold value, said outputs produced, respectively, by a last said serial constituent decoder (102) performing a first and a subsequent iteration and a first said serial constituent decoder (100) performing said subsequent iteration; and

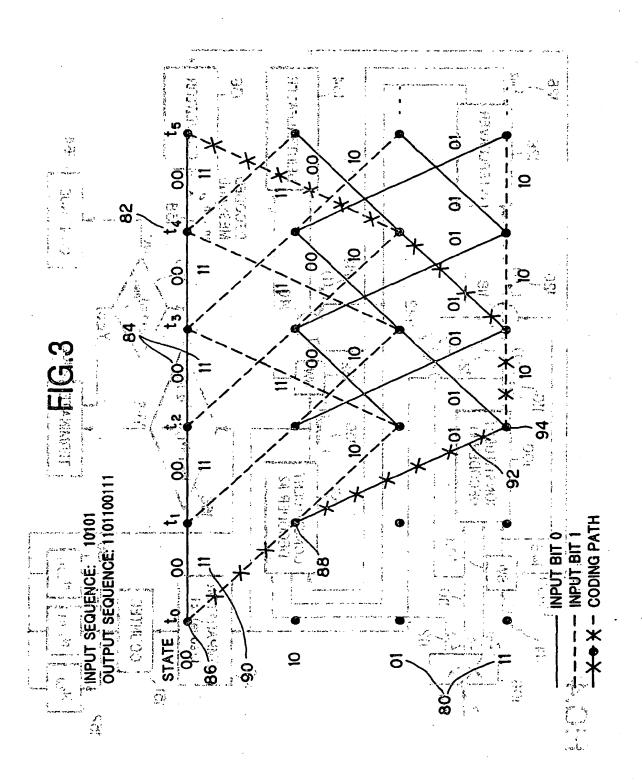
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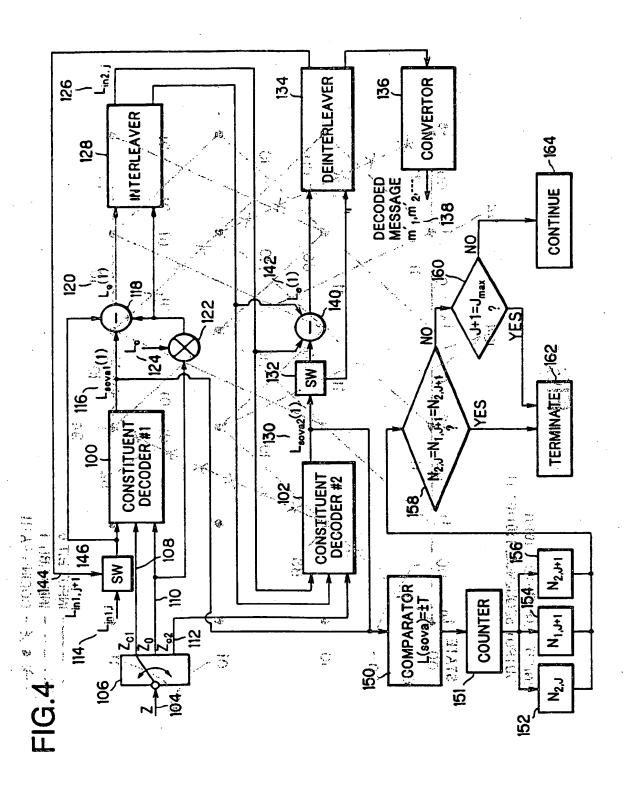
(e) a decision unit (158) to terminate the operation of said turbo code decoder when said first, second, and third numbers of said outputs are approximately equal to the control of the c

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EUROPEAN SEARCH REPORT

Application Number EP 00 10 1174

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	PROCEEDINGS 1998 IE SYMPOSIUM ON INFORM 16 = 21 August 199	ATION THEORY ,	5987	
	CAMBRIDGE, MA, USA * page 279, left-ha paragraph - right-h	nd column, last		(х.д» - ¹
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	* page 1188, right- line 30 *	hand column, line 2	23° -	
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X : par Y : par	ATEGORY OF CITED DOCUMENTS ticularly relevant if taken alone ticularly relevant if combined with anotument of the same category	E : earlier after the ther D : docume	or principle undorlying the patent document, but put of filing date ant cited in the application tried for other reason	blished on, or on

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European Patent EUROPEAN SEARCH REPORT

Application Number EP 00 10 1174

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ANNEX TO THE EUROPEAN SEARCH REPORT ON EUROPEAN PATENT APPLICATION NO.

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This annex lists the patent family members relating to the patent documents cled in the above-mentioned European search report. The members are as contained in the European Patent Office EDP the only a Sufficient of the purpose of information.

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EP 00 10 1174

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